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Jurnal Teknik Informatika dan Sistem Informasi

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Rizki Rahmat Rihsyah, Youllia Indrawaty Nurhasanah, Irma Amelia Dewi

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Performance Analysis of NFS Protocol Usage on VMware ESXi Datastore

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Abstract — Hypervisor virtualization that uses bare metal architecture allows to allocate and provide resources for each created virtual machines. Resources such as: CPU and memory, can be added or upgraded anytime to the host hardware (virtualization server) to be able to create more virtual machines. However, upgrading the hard drive size cannot be done anytime if there are already have data or virtual machine that has fully operated on the host hardware, related to the raid system and the establishment of hard drive partition.

Upgrading hard drive size on virtualization server can be done by using NFS protocol on NAS server. VSphere ESXi able to use NFS protocol and store the virtual disk that is used by virtual machine as guest operating system on network storage besides using local hard drive (host hardware hard drive). When the virtual machine want to run the guest operating system, it will request to write/read virtual disk there is stored on NAS by using NFS protocol through the network.

In this research, measurements has been taken on data communication performance due the usage of NFS as virtual machine's datastore in addition to local hard drive usage on server's device. Measurements were performed by sending various data size from client to server (virtual machine) and measure write/read speed on server's cache memory and hard drive, and measure RTT (Round-trip Time) delay between client-server. The testing has been conducted on virtual machines that use local drive and NFS as virtual disk datastore.

Keywords— datastore, ESXi, NAS (Network Attached Storage), network , NFS (Network File System), performance, virtualization, vSphere.

I. INTRODUCTION

The virtualization concept already widely used for server deployment. One machine can create multiple heterogenous operating system depends on needed. The using of virtualized environment makes deployment more efficient such less power cost then using unvirtualized environment [1]. However, virtualization requires more resources such CPU, memory, and storage space. Each creation of virtual machine will be given CPU, memory, storage, network access, from available resources which provided by the hypervisor. The resources are guaranteed only can be used for concerned virtual machine by the hypervisor [2] [3]. Since CPU, memory, and network adapter are depends on the type of the

server, it makes upgrading them are limited. Storage is also limited in the number of hard drive's slot.

Many storage externalization technology such as NAS (Network Attached Network) and SAN (Storage Attached Network) using NFS protocol already widely used [4]. NFS frequently used to allow user on a client computer to access file through the network like using local storage [5]. The use of external storage is possible to use such as NAS on vSphere datastore to provide more storage space for virtual machines.

NFS implementation as virtual machine datastore can improve flexibility in the process of further development. However, since every data of virtual machine will be stored at different location, it will degrade the output performance such as data communication speed performance. Our research will measure data transfer performance by using NFS as vSphere ESXi datastore such as write/read speed to virtual machine's cache memory and hard drive, and measure RTT (round-trip time delay) from client side to virtual machine. Testing will be performed on the virtual machines that use datastore on direct attached hard drive and NFS protocol.

II. RELATED WORK

Many studies related to the field of virtualization performance and efficiency have been conducted [1] [6]. Maria and Hammad made a research by deployment NAS on ESXi virtual machine uses virtual switch on it. Their research focus on the procedures of NAS implementation on virtual environment. They create a virtual machine as NAS and configure NFS protocol for the procedures for accessing the NAS [4]. In our research, we are not create NAS on ESXi but use NAS as ESXi virtual machine datastore by using NFS for write/read communication protocol.

Debabrata and Rajesh provide a model on their research to make an encrypted virtual machine storage on the scope of the private cloud. They realize the importance of using virtual machine to reduce cost during maintenance process and expanding the infrastructure [7]. Our research focused on storage addition efficiency, different from their research that provide storage security model.

Hasan, Lue, and Martin perform benchmarking on the 3 types of virtual machine (Hyper-V, VMware ESXi, and Xen Hypervisor). They did a benchmark on the virtualization approach, architecture, hardware, and software on each

virtual machines. They was measuring virtual machines performance based on thread clock tick and memory load. In our research, benchmarking on virtual machine focused on measuring write/read speed on cache memory and hard drive before and after using NFS as virtual machine datastore on ESXi [8].

Vaughn, Michael, Larry, and Peter recommended many methods on their report for improving vSphere storage including using NFS as datastore. However, they aren't test the impact by using NSF through the value of data communication performance [5].

III. LITERATUR REVIEW

A. Virtualization

Since 1960, virtualization has been used for dividing the system resources provided by mainframes computer between different applications [9] [10]. Virtualization is considered to be act of abstracting the physical boundaries of a technology [4]. Virtualization has a major impact on the data center by decreasing a number of physical servers need to run on back office and greatly increasing the manageability and flexibility of that infrastructure. Since there are occurrence of physical abstraction such improving security, flexibility and performance, makes one virtual machine separate from one other [7]. While one of them fails, it will completely isolated from all other on a physical machine, including other virtual machines. This is make the problem will be contained and increase the security and manageability [11]. There are several approach of virtualization based on how the virtual machines are controlled [7].

1) Operating System-Based Virtualization

In Figure 1 shows the illustrated architecture of operating system-based virtualization. Virtualization is provided by host operating system. It support feature to isolated virtualized guest OS's on a single physical host hardware even using the same operating system kernel and characteristics. The virtual machines are controlled by the host operating system.

The using of operating system-based virtualization have several weaknesses such performance degradation. Since guest operating system need to performs I/O operation, it need to be interpreted by the host operating system before accesing the hardware which can access the kernel. Extra CPU is needed and lead to performance degradation. Operating system-based virtualization is not flexible as other virtualization approaches since it can't host a quest OS different with the host one, or a different guest kernel [12].

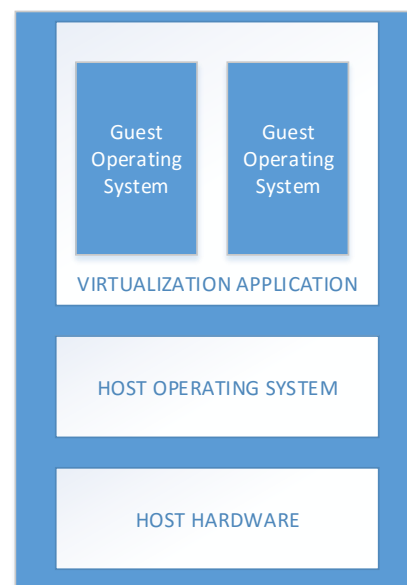


Fig 1. Architecture of operating system-based virtualization [1]

2) Application-Based Virtualization

Architecture of application-based virtualization is shown in Figure 2. The computer programs are encapsulated in a software underlying operating system on which it is executed. Fully application virtualization requires a virtualization layer. Application virtualization layer replace part of the runtime environment normally provided by the operating system. The layer intercepts all disk operations of virtualized applications and tranparently redirects them to a virtualized location [13].

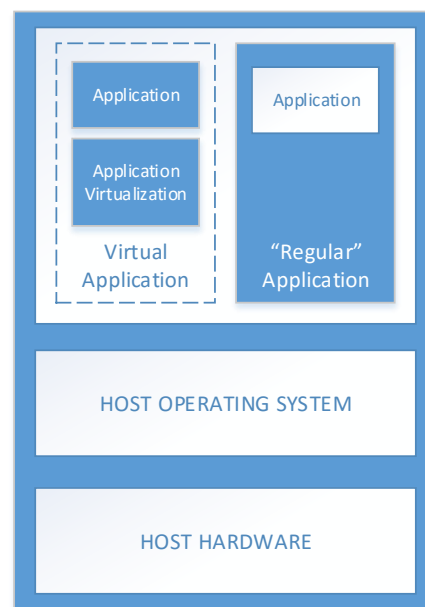


Fig 2. Architecture of application-based virtualization [1]

Application virtualization helps solve the problem of application incompatibility in an operating system and others by adding a layer of isolation between operating system and the application. This virtualization approach gives administrators a way to create virtualized applications that help minimized risk, save time, increase user satisfaction, and reduce cost for support and regression testing [14].

There are several limitation of application virtualization [12]:

- i. Not all software can be virtualized.
- ii. Only file and registry-level compatibility issues between legacy applications and newer operating system can be addressed by application virtualization.
- iii. Licensing application virtualization must be correctly licensed between virtualized software and virtualized applications.

3) Hypervisor-Based Virtualization

One of many hardware virtualization such hypervisor allow multiple operating systems, termed guest, to run concurrently on a host computer. The hypervisor present to the guest operating system a virtual operating platform and monitors the execution of the guest operating systems [7]. The guest software is not limited to user applications, many host allow the execution of complete operating systems. The guest software executes as if it were running directly on the physical hardware [12].

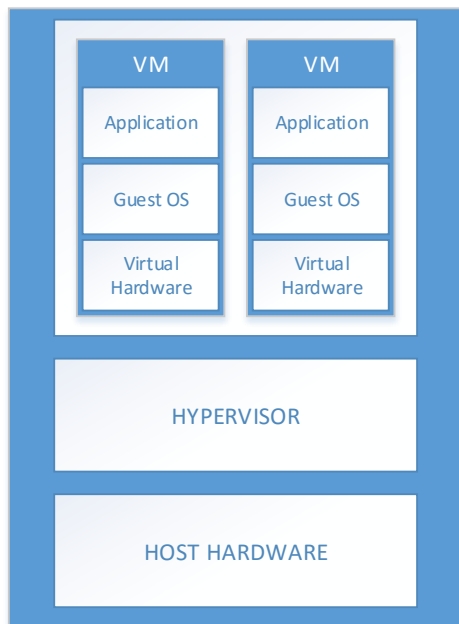


Fig 3. Architecture of bare metal hypervisor-based virtualization [1]

Hypervisor or Virtual Machine Monitor (VMM) is installed on server hardware whose only task to run guest operating systems. The virtualization layer directly control the hardware and manages guest operating systems. Since bare metal model (Figure 3) directly implements the virtualization in the hardware level, made the system overhead transferring I/O operation from guest operating system to the hardware is less than hosted model (Figure 4).

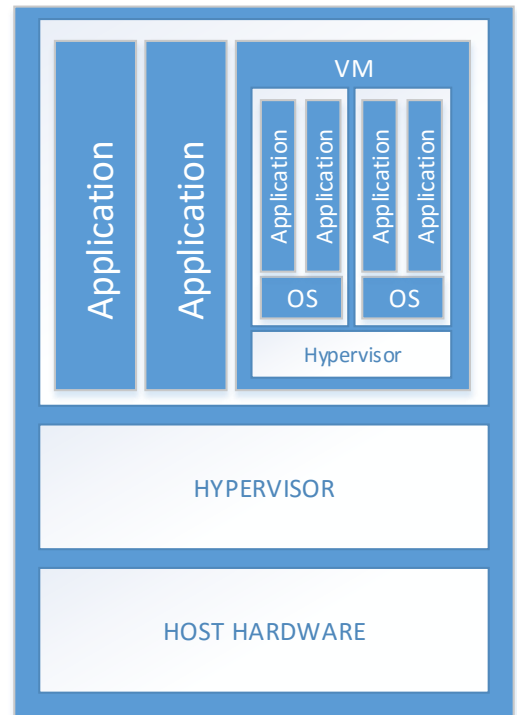


Fig 4. Architecture of hypervisor-based virtualization [1]

The bare metal hypervisors are installed directly on host hardware, just like any other operating system [15]. The size of management software is very small, it makes the resource used by the virtualization layer can be ignored. Since the virtual machines are not built within a host operating system, it is more flexible and reliable for various applications. The hypervisor is available at the boot time of machine in order to control the sharing of system resources across multiple virtual machines [7].

B. Advantage of Virtualization

1) Cost

Hardware is the highest cost when we build a system. While we can reduce the amount of hardware, that means reduce our cost. It is possible to achieve reductions by consolidation smaller servers into more powerful servers. Cost reduction obtained from

hardware cost reductions, operation cost reductions, floor space, and software licence. The use virtualization can reduce overall cost from 29 to 51%. [1] [7] [16]

2) *Reliability*

A software or operating system failure in a virtual machine does not affect other virtual machines [7].

3) *Security*

The code for hypervisor and boot operating system kept as small as possible, creates a smaller attack surface. Virtual machines run at lower level of permissions than hypervisor, it makes inhibits virtual machines attacks [17].

4) *Load Balancing*

The software state of entire virtual machines is completely encapsulated by the virtual machine management (VMM), it is made relative easy to migrate virtual machine to ether platform in order to improve performance through better load balancing [18]. Since the hypervisor can provide hardware driver from its vendor, the using of multiple network interface will automatically trigger load balancing without any configuration by ESXi it self.

C. *vSphere ESXi*

VMWare ESXi is the next-generation hypervisor, providing a new foundation for virtual infrastrcuture. This innovative architecture operates independently from any general purpose operating system, offering improved security, increased realibility, and simplyfied management [19].

Server virtualization using VMWare's ESXi is the most dominant and stable virtualization technology. VMWare alone holds 50% of the market share in virtualization, whereas other vendors combine the remaining 50%. A research survey done by F5 networks in the time from September to Desember 2008 reported that VMWare ESXi is the most widely deployed server virtualization product [20].

Besides being the most dominant vendor in the virtualization market, VMWare's ESXi have a features that make it reliable and favourite choice for IT professionals [15]:

i. *Small foot print*

ESXi just need 70MB of disk space for instalation. Compared with other bare metal model such Hyper-V with minimum 2GB of disk space and 1.8GB with XenServer v5.6 [8].

ii. *Hardened Drivers*

Hardware drivers for virtual machine already optimized by hardware vendors.

iii. *Advance memory management*

Feature to reclaim unused memory, duplicate memory pages, and compress memory pages.

iv. *Advance storage management*

Feature to provides independent storage management. The allocation of storage can be done without interrupting the other virtual machines or host machine.

v. *Host resource management*

Feature for network traffic shaping, resource sharing for each virtual machines, and allows the setting of quality of service priorities for storage and network I/O.

vi. *Flexible resource allocation*

The resource allocation can be done on the fly, such as add virtual CPU, memory, and virtual disk (hard disk space)

ESXi install and runs without the Linux-based Service Console like VMWare ESX. This makes ESXi an ultra light footprint of approximately 70MB. ESXi provide all the same virtualization feature that VMWare ESX provided in earlier versions. The reason that ESXi is able to support the same extensive set of virtualization as VMWare ESX without the Service Console is that the core of virtualization functionality was not found in the Service Console [8]. VMKernel manages the virtual machines access to the underlying physical hardware by providing CPU scheduling, memory management, and virtual switch data processing [21]. Figure 5 shows the architecture of next generation VMWare virtualization foundation ESXi after ESX.

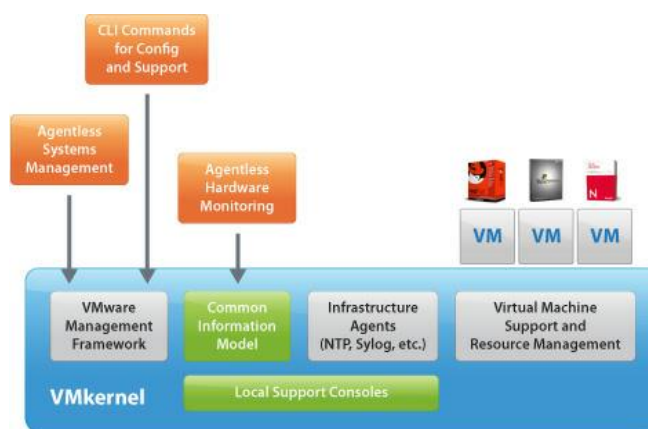


Fig 5. Architecture of VMWare ESXi [16]

D. *NAS*

Network attached storage (NAS) is a term used to refer to storage elements or devices that connect to a network and provide file access services to computer systems. NAS devices attach directly to networks, such as LAN (local area network), by using TCP/IP protocol and serve files to any client connected to the network. A client computer access NAS devices uses a file system device driver to access the stored data. The file system device driver typically uses a file access protocol such as Common Internet File System (CIFS)

or Network File System (NFS). NAS devices interpret these command and perform the internal file and device I/O operation necessary to execute them [22][23].

E. NFS

The Network File System (NFS) is designed to be portable across different machines, operating systems, network architectures and transport protocol to provides remote access to shared files across networks. This portability is achieved through the use or Remote Procedure Call (RPC) primitives built on top of eXternal Data Representation (XDR) [24].

Since the NFS protocol was intended to be a stateless connection, that is made the server should not need to maintain any protocol state information about any of it’s clients in order to function correctly. Stateless servers have a distinct advantage over stateful servers in the event of failure. A client just need to retry a request until the server responds, it does not even need to know that the server has crashed, or the network connection went down. The client of a stateful server needs to either detect a server failure and rebuild the server’s state when it comes back up, or cause operations to fail [24].

All of the procedures in the NFS protocol assumend to be synchronous. When a procedure returns to the client, the client can assume that the operation has completed and any data associated with the request is now on stable storage.

IV. EXPERIMENTAL SETUP

A. Hardware Configuration

The hardware platform used for conduction this research has the following spesification.

- i. ESXi (Virtualization Server): HP ProLiant DL380 Gen9, Intel® Xeon® CPU E5-2650v3 (20 CPUs, 40 Logical Processor) @ 2.30GHz with 25MB SmartCache, 64GB memory, 4TB hard drive, 6 ethernet port with spec: 4 Broadcom NetXtreme Gigabit Ethernet Port and 2 Broadcom QLogic 57810 10 Gigabit Ethernet Port.
- ii. NAS Server: HP StoreEasy 1550, Intel® Xeon® CPU E5-2603v3 (6 CPUs) @ 1.6GHz, 8GB Memory, 16TB hard drive, 2 Gigabit Ethernet Broadcom BCM5715 NetXTreme.

B. Software Configuration

Hypervisor provided by using lisenced VMware vSphere 6.0 (unlimited cores per CPU). These version where the last shipping releases at the time of doing this research (started in January 2017).

C. Guest Operating System

Windows Server 2012 (64bit) is the guest operating system used on our research. There are two Windows Server 2012, first is stored at virtualization server datastore directly by using local hard drive in it, the second is stored at NAS

server by configuring NFS protocol between virtualization server and NAS server for write/read communication protocol.

Windows Storage Server 2012 (64bit) R2 installed at NAS server which provide storage for any client by using NFS protocol.

Windows 10 pro (64bit) installed at Client PC and used for communication testing to both Windows Server 2012.

D. Virtual Machine Configuration

Since this research using ESXi as hypervisor, the created virtual machine will be fully-virtualized instead of paravirtualized. Each virtual machine is configured to have one virtual CPU, 4GB of memory, 40GB hard drive for I/O file testing, and 1 Gigabit ethernet port.

Datastores

Identification	Device	Type	Capacity
Local Drive	HP Serial Attached SCSI Disk (naa.60...	VMFSS	2.18 TB
NASStorage	10.10.10.30.60:/tes	NFS	10.68 T

Fig 6. ESXi datastore (VMFSS and NFS)

Figure 6 shows the ESXi have 2 types of datastore. First virtual machine will be installed at “Local Drive”, and the other one will be installed at “NAS Storage”. The communication between ESXi and NAS server is using NFS protocol that allowed ESXi to access shared directory and store virtualization image (.vmdk, .vmx, .nvram, etc) on NAS.

E. Network Topology

The hypervisor installed at server device and stored at DMZ area and NAS devices as ESXi datastore stored at local area network. Figure 7 shows the connection between ESXi, NAS devices, and PC Client.

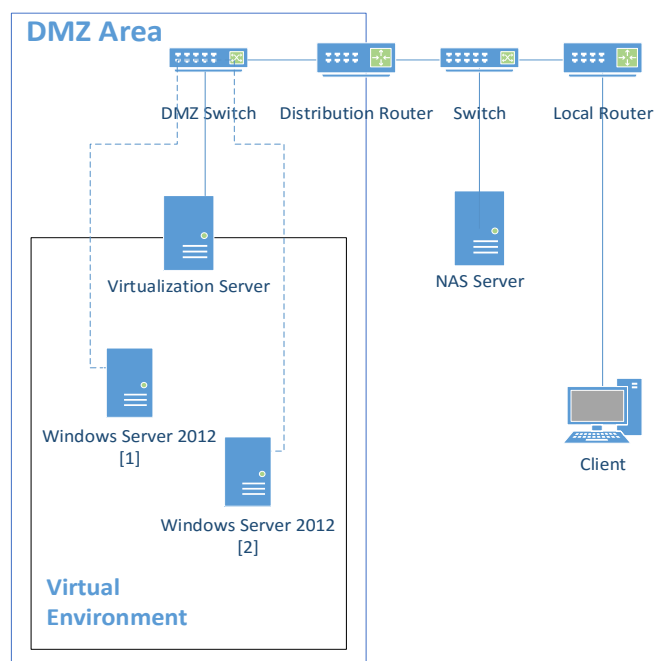


Fig 7. Network Topology

Even the both Windows Server 2012 are running at virtual environment under ESXi hypervisor, one of it (Windows Server 2012 [2]) is stored at NAS server besides on virtualization server's local drive. Every times Windows Server 2012 [2] runs, the virtualization server will be ask the files needed (virtual disk) to NAS server by using NFS protocol. Figure 8 shows the NFS protocol communication between virtualization server to NAS server, captured by using Wireshark software.

```

122.87 30.60 NFS 1514 V3 WRITE Call (Reply In 912)
122.87 30.60 NFS 1514 V3 WRITE Call (Reply In 977)
30.60 122.87 NFS 230 V3 WRITE Reply (Call In 616)
122.87 30.60 NFS 1514 V3 WRITE Call (Reply In 1040)
30.60 122.87 NFS 230 V3 WRITE Reply (Call In 647)
    
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Fig 8. Communication between virtualization server and NAS server

V. TESTING PROCEDURES AND RESULTS

Windows Server 2012 is selected as guest operating system since we used a software that used TCP/IP protocol for transmitting and receiving data between client-server and also measuring write/read data communication speed.

A. Testing Metrics

There are two mechanism of data communication on this test. Write and read data are performed on server's cache memory and hard drive. Measurements are performed at write/read speed and RTT (Round-trip Time) between client-server.

B. Measuring Process

The measurement on each testing metrics are performed on both virtual machines to obtain the performance difference between using local hard drive and NFS protocol as virtual machines datastore. The measurement will be conducted 1000 times for each scenarios. We conducts test by using four data sizes (100KB, 1M, 10M, 100M) send over network between client-server. For each test, we generate 100 data packets on every data size and do 10 times testing.

C. Testing Scenarios

i. Write/Read from Client to Server's Cache Memory

We have been conduct several testing scenarios by using four data sizes. It intended to measure the virtual machine's write/read capabilities to it's cache memory and hard drive. We measure the write/read performance to virtual machine that use local drive at host machine and NFS protocol as virtual machine's datastore.

The first testing scenario is sending data from client (PC Client) to server (Windows Server 2012 [1] & [2]) over the network by using TCP/IP protocol. Figure 9 to 16 show the performance different on write time (ms) to server's cache memory on 100KB data size on virtual machine which use local hard drive (VMFSS) and NFS protocol as virtual machine datastore to storing virtual disk (.vmdk).

The Figure 9 shows the write process to server's cache memory using local hard drive is more unstable compared than using NFS protocol.

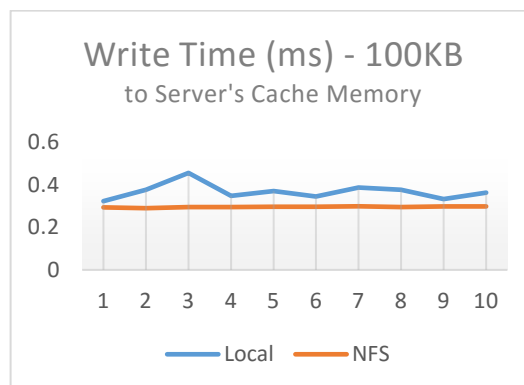


Fig 9. Write Time to Cache Memory – 100KB data size

The average of write time (ms) from client to server with 100KB data size is shown in Table I. Every rows shows the average of 100 times testing. Average of 1,000 times write testing with 100KB data size to server's cache memory shows that using NFS protocol as virtual machine datastore is 24.38% faster than using local drive at host hardware.

TABLE I
WRITE TIME TO SERVER'S CACHE MEMORY - 100KB DATA SIZE

No.	Local Drive	NFS
1	0.3227733	0.2932900
2	0.3760763	0.2896637
3	0.4554326	0.2945069
4	0.3478442	0.2956532
5	0.3703460	0.2957401
6	0.3447231	0.2963665
7	0.3861392	0.2986223
8	0.3761321	0.2956834
9	0.3331197	0.2974705
10	0.3626582	0.2978439
Avg	0.36752447	0.29548405

Figure 10 shows the result of read time (ms) from client to server's cache memory on 100KB data size. The graph shows that read time by using NFS protocol more stable compared by using local hard drive on host hardware.

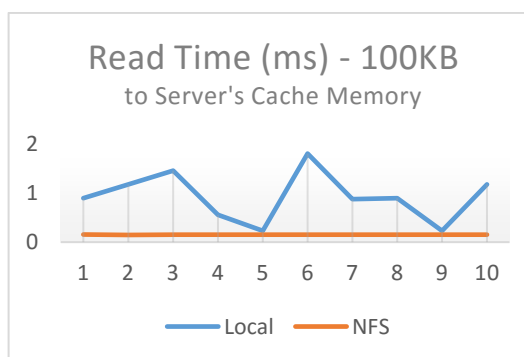


Fig 10. Read Time to Cache Memory – 100KB data size

The average of read time (ms) needed from client to server's cache memory with 100KB data size is shown in Table II. Average of 1000 times read testing with 100KB data size to server's cache memory shows that using NFS as virtual machine datastore is 509.67% faster than using local drive at host hardware.

TABLE II
READ TIME TO SERVER'S CACHE MEMORY - 100KB DATA SIZE

No.	Local Drive	NFS
1	0.8973481	0.1562289
2	1.1820205	0.1465906
3	1.4616412	0.1541129
4	0.5558776	0.1507121
5	0.2318107	0.1543894
6	1.8124936	0.1529284
7	0.8802639	0.1535888
8	0.9002681	0.1537114
9	0.2330693	0.1541873
10	1.1792473	0.1545411
Avg	0.93340403	0.15309909

Figure 11 shows the value of write time (ms) from client to server's cache memory on 1MB data size. The graph shows the using of local hard drive are more stable compared with NFS protocol on 1MB than 100KB data size.

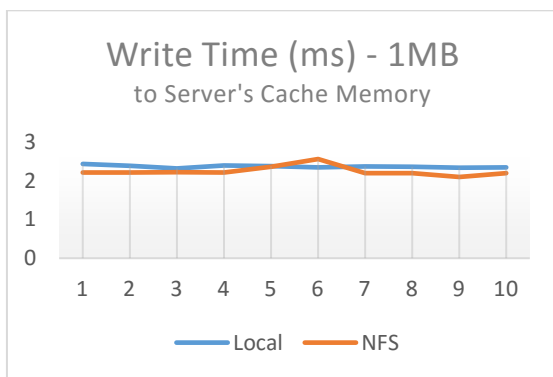


Fig 11. Write Time to Cache Memory – 1MB data size

The average of write time (ms) needed from client to server with 1MB data size is shown in Table III.

TABLE III
WRITE TIME TO SERVER'S CACHE MEMORY – 1MB DATA SIZE

No.	Local Drive	NFS
1	2.4297315	2.2070025
2	2.3809750	2.2113293
3	2.3138561	2.2205531
4	2.3894056	2.2096884
5	2.3703614	2.3598893
6	2.3432618	2.5560472
7	2.3619025	2.1904816
8	2.3555105	2.1920535
9	2.3333875	2.0934884
10	2.3435219	2.1927912
Avg	2.36219138	2.24333245

Average of 1000 times write testing with 1MB data size to server's cache memory shows that using NFS as virtual machines datastore is 5.29% faster than using local drive at host hardware.

Figure 12 shows the value of read time (ms) from client to server's cache memory on 1MB data size. The graph shows the read process on 1MB data size by using NFS protocol are more stable and faster than using local drive as virtual machine datastore.

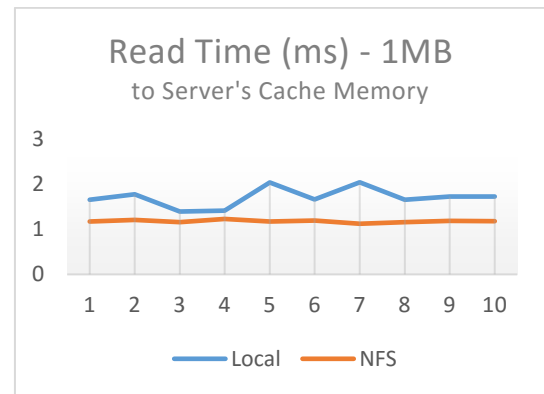


Fig 12. Read Time to Cache Memory – 1MB data size

Table IV shows the average of 1000 times read testing on 1MB data size to server's cache memory. The using of NFS is 45.01% faster than using local drive as virtual machine datastore.

TABLE IV
READ TIME TO SERVER'S CACHE MEMORY – 1MB DATA SIZE

No.	Local Drive	NFS
1	1.6581289	1.1766468
2	1.7769015	1.2073970
3	1.3960516	1.1613709
4	1.4138687	1.2313505
5	2.0455276	1.1735411
6	1.6627443	1.1972251
7	2.0456758	1.1247076
8	1.6585109	1.1625198
9	1.7281476	1.1858430
10	1.7276530	1.1800396
Avg	1.71132099	1.18006414

Figure 13 shows the value of write time (ms) from client to server on 10MB data size. The graph shows the write testing with 10MB data size to server's cache memory on local drive and NFS are both unstable value compared with previous test.

The Table V shows the average of 1000 times write testing on both servers. The using of local drive is 2.12% faster than using NFS as virtual machine datastore. Compared with previous testing on smaller datasize, write data to server's cache memory that using local drive as virtual machine datastore on 10MB data size show faster results than using NFS protocol.

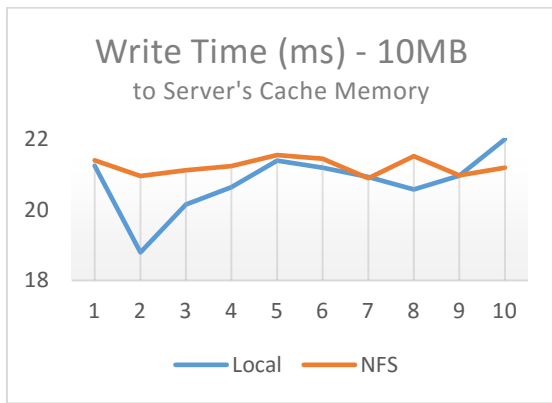


Fig 13. Write Time to Cache Memory – 10MB data size

TABLE V
WRITE TIME TO SERVER'S CACHE MEMORY – 10MB DATA SIZE

No.	Local Drive	NFS
1	21.2407821	21.4034873
2	18.7942882	20.9534845
3	20.1486679	21.1149016
4	20.6350253	21.2316414
5	21.3881703	21.5485224
6	21.1889365	21.4456097
7	20.9265949	20.8925960
8	20.5718300	21.5166040
9	20.9605279	20.9692456
10	21.9976195	21.1947304
Avg	20.78524426	21.22708229

Figure 14 shows the value of read time (ms) from client to server 10MB data size. The graph shows the using local drive are more unstable than NFS as datastore.

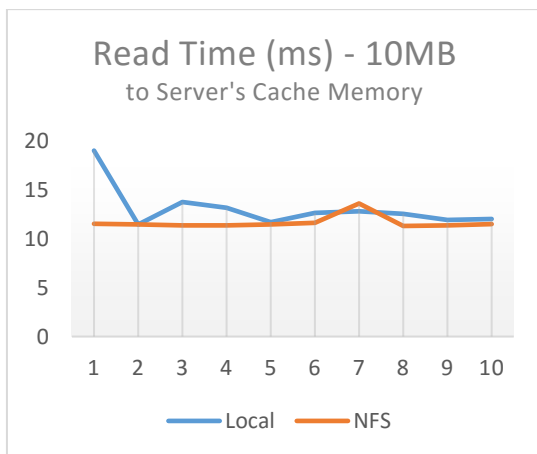


Fig 14. Read Time to Cache Memory – 10MB data size

The Table IV shows the result of 1000 times testing read data from client to server's cache memory on 10MB data size by using of NFS as virtual machine datastore is 12.35% faster than using local drive.

TABLE VI
READ TIME TO SERVER'S CACHE MEMORY – 10MB DATA SIZE

No.	Local Drive	NFS
1	18.9812338	11.5216015
2	11.4377432	11.4557662
3	13.7364794	11.3425241
4	13.1597428	11.3561308
5	11.6862156	11.4530614
6	12.6336952	11.6170974
7	12.7788510	13.5841254
8	12.5122481	11.2879391
9	11.9090408	11.3533452
10	12.0038139	11.4834139
Avg	13.08390638	11.6455005

Figure 15 shows the value of write time (ms) from client to server's cache memory on 100MB data size. The graph shows that the using of local drive and NFS are both unstable than smaller data size on previous test.

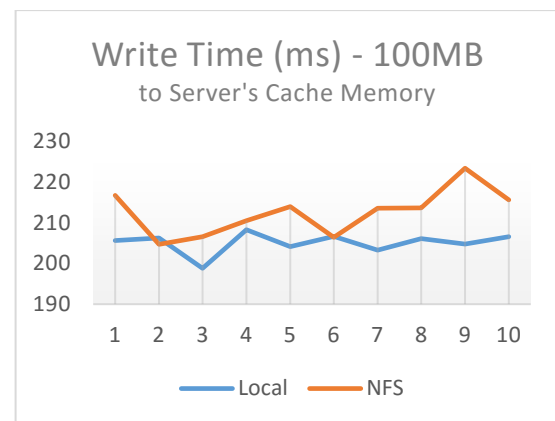


Fig 15. Write Time to Cache Memory – 100MB data size

The Table VII shows the average of 1000 times write testing to server's cache memory on 100MB data size by using local drive as virtual machine datastore is 3.51% faster than NFS.

TABLE VII
WRITE TIME TO SERVER'S CACHE MEMORY – 100MB DATA SIZE

No.	Local Drive	NFS
1	205.6173631	216.7695078
2	206.3029872	204.7207051
3	198.8206734	206.6124627
4	208.3145364	210.5396421
5	204.1249096	213.9944990
6	206.6413029	206.4118865
7	203.3173913	213.6243310
8	206.1648641	213.6666022
9	204.7953219	223.4595259
10	206.6063254	215.6265829
Avg	205.0705675	212.5425745

Figure 16 shows the values of read time (ms) from client to server's cache memory on 100MB data size. The graph clearly shows that the using local drive is need less time than NFS on read 100MB data size.

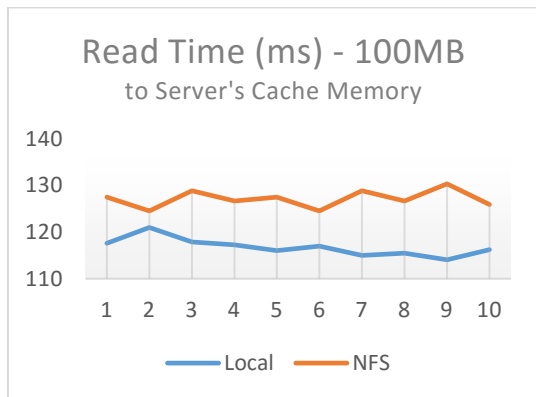


Fig 16. Read Time to Cache Memory – 100MB data size

The Table VIII shows the average of 1000 times read testing on 100MB data size to server's cache memory. The test results show that the using of local drive as virtual machine datastore is 8.90% faster than using NFS on 100MB read test scenario.

TABLE VIII
READ TIME TO SERVER'S CACHE MEMORY – 100MB DATA SIZE

No.	Local Drive	NFS
1	117.6023214	127.4876240
2	120.9903297	124.5415373
3	117.8610673	128.9027952
4	117.2513944	126.6516907
5	116.0152739	127.4876240
6	116.9644377	124.5415373
7	114.9893922	128.9027952
8	115.4930069	126.6516907
9	114.0547803	130.3600002
10	116.2309486	125.9291211
Avg	116.7452952	127.1456416

ii. Write/Read from Client to Server's Hard Drive

The second testing scenario is sending data from client to server's hard drive. Figure 17 to 24 shows the performance different on write/read time (ms) from client to server which using local datastore (VMFSS) and NFS protocol as virtual machine datastore.

The Figure 17 shows the results of 1000 times write testing on 100KB data size to server's hard drive. Same as previous scenario, for every 100 times testing will be conducted average calculation. The graph clearly shows that the using local drive is faster than NFS as virtual machine datastore on write 100KB data size to server's hard drive.

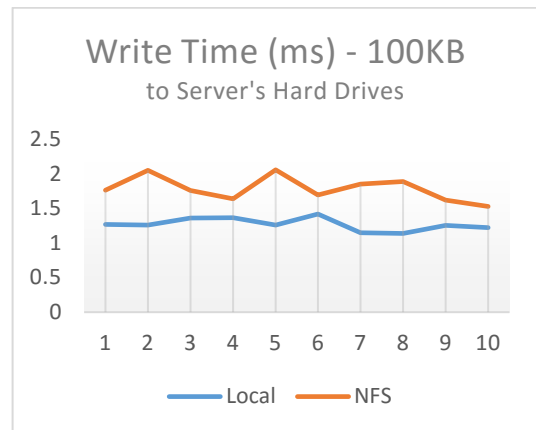


Fig 17. Write Time to Server's Hard Drive– 100KB data size

The Table IX shows the average of 1000 times write testing results on server's hard drive on 100KB data size. The results show that the write time on virtual machine that used local drive as virtual machine datastore is 40.58% faster than using NFS.

TABLE IX
WRITE TIME TO SERVER'S HARD DRIVE – 100KB DATA SIZE

No.	Local Drive	NFS
1	1.2692339	1.7658046
2	1.2613508	2.0487450
3	1.3590300	1.7592541
4	1.3642295	1.6378762
5	1.2587541	2.0578476
6	1.4187532	1.6919528
7	1.1473697	1.8511347
8	1.1379583	1.8861857
9	1.2552401	1.6193669
10	1.2229089	1.5283764
Avg	1.26948285	1.7846544

The Figure 18 shows the 1000 times read testing on 100KB data size on server's hard drive. The graph shows the read process on virtual machine that use NFS as datastore is unstable and slower compared with using local drive.

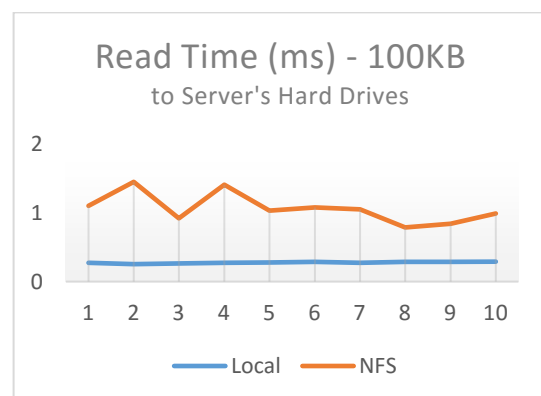


Fig 18. Read Time to Server's Hard Drive – 100KB data size

The Table X shows the average of 1000 times read testing to server's hard drives. The result shows that using local drive as virtual machine datastore is 285.12% faster than using NFS.

TABLE X
READ TIME TO SERVER'S HARD DRIVE – 100KB DATA SIZE

No.	Local Drive	NFS
1	0.2727091	1.1009283
2	0.2532925	1.4479928
3	0.2650531	0.9169837
4	0.2740736	1.4086522
5	0.2796387	1.0301360
6	0.2850377	1.0751110
7	0.2731174	1.0475101
8	0.2846108	0.7864894
9	0.2864403	0.8413505
10	0.2896118	0.9882189
Avg	0.2763585	1.06433729

Figure 19 shows the values of write time (ms) from client to server's hard drives on 1MB data size. The graph clearly shows that the using local drive is need less time than NFS on writing 1MB data size to server's hard drives.

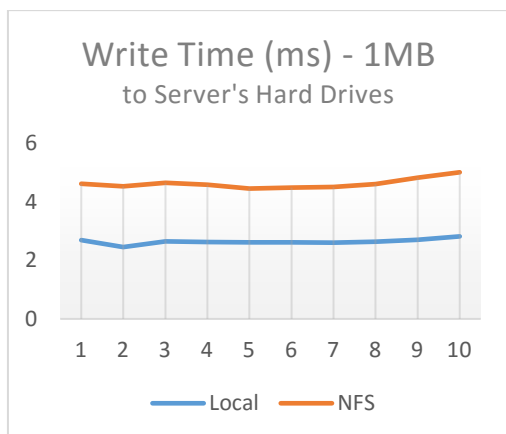


Fig 19. Write Time to Server's Hard Drive – 1MB data size

TABLE XI
WRITE TIME TO SERVER'S HARD DRIVE – 1MB DATA SIZE

No.	Local Drive	NFS
1	2.6951906	4.6125644
2	2.4538437	4.5276795
3	2.6429746	4.6541678
4	2.6237479	4.5839958
5	2.6098378	4.4526557
6	2.6108253	4.4808464
7	2.6094093	4.5075496
8	2.6389205	4.6009523
9	2.7065288	4.8241163
10	2.8184016	5.0088768
Avg	2.64096801	4.62534046

The Table XI shows the result of measuring write speed from client to server's hard drive on 1MB data size. The test conducted 1000 times and shows that the write time (ms)

using local drive as virtual machine datastore is 75.13% faster than NFS on writing 1MB data size.

Figure 20 shows the values of read time (ms) from client to server's hard drives on 1MB data size. The graph clearly shows that the using local drive is need less time than NFS on reading 1MB data size.

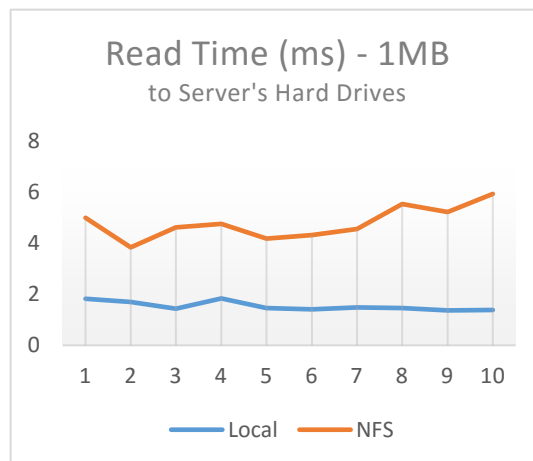


Fig 20. Read Time to Server's Hard Drive – 1MB data size

The Table XII shows the result of measuring read speed from client to server's hard drive on 1MB data size. The test conducted 1000 times and shows that the read time (ms) using local drive as virtual machine datastore is 212.13% faster than NFS on reading 1MB data size.

TABLE XII
READ TIME TO SERVER'S HARD DRIVE – 1MB DATA SIZE

No.	Local Drive	NFS
1	1.8280787	4.9922166
2	1.7018325	3.8379347
3	1.4340294	4.6196290
4	1.8379318	4.7546660
5	1.4574793	4.1798378
6	1.4069189	4.3161234
7	1.4892594	4.5611101
8	1.4590108	5.5375688
9	1.3681861	5.2294828
10	1.3829000	5.9326323
Avg	1.53656269	4.79612015

Figure 21 shows the values of write time (ms) from client to server's hard drives on 10MB data size. The graph clearly shows that the using local drive is need less time than NFS on writing 10MB data size.

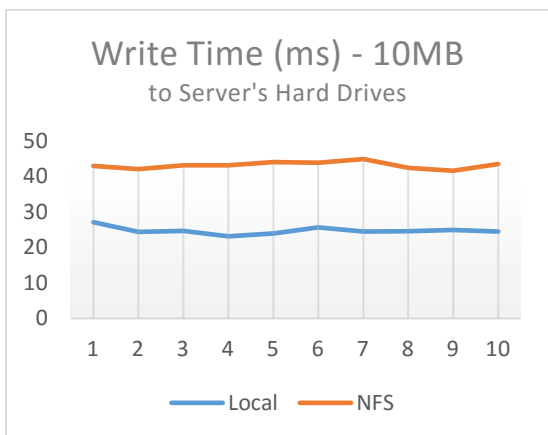


Fig 21. Write Time to Server's Hard Drive – 10MB data size

The Table XIII shows the result of measuring write speed from client to server's hard drive on 10MB data size. The test conducted 1000 times and shows that the write time (ms) using local drive as virtual machine datastore is 74.52% faster than NFS on writing 10MB data size.

TABLE XIII
WRITE TIME TO SERVER'S HARD DRIVE – 10MB DATA SIZE

No.	Local Drive	NFS
1	27.1478046	43.0757329
2	24.4528828	42.1070175
3	24.6816513	43.2469012
4	23.1744430	43.2497131
5	24.0107874	44.0722654
6	25.6908536	43.9439595
7	24.5191706	44.9724713
8	24.6346203	42.5227648
9	24.9806222	41.6700380
10	24.4733174	43.5520640
Avg	24.77661532	43.24129277

Figure 22 shows the values of read time (ms) from client to server's hard drives on 10MB data size. Same as the previous test with smaller data size, the graph clearly shows that the using local drive is need less time than NFS on writing 10MB data size.

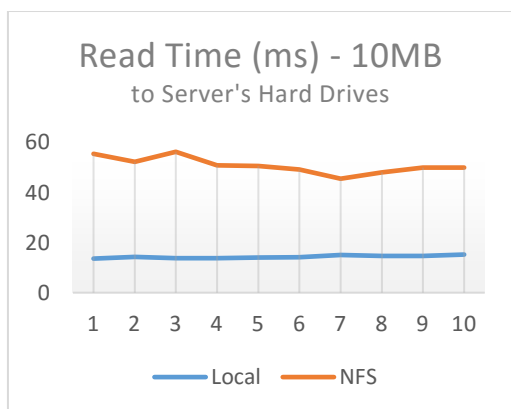


Fig 22. Read Time to Server's Hard Drive – 10MB data size

TABLE XIV
READ TIME TO SERVER'S HARD DRIVE – 10MB DATA SIZE

No.	Local Drive	NFS
1	13.6111554	55.4404134
2	14.2858530	52.3000522
3	13.7428692	56.3048110
4	13.7763704	50.8230918
5	14.0682058	50.6481013
6	14.1392855	49.2817658
7	15.0144403	45.5446373
8	14.6455623	48.1021028
9	14.6736816	49.9587382
10	15.2351220	49.9417438
Avg	14.31925455	50.83454576

The Table XIV shows the average of 1000 times read testing on 10MB data size to server's hard drives. The test results show that the reading time using of local drive as virtual machine datastore is 255.00% faster than using NFS on 10MB data size.

Figure 23 shows the values of write time (ms) from client to server's hard drives on 100MB data size.

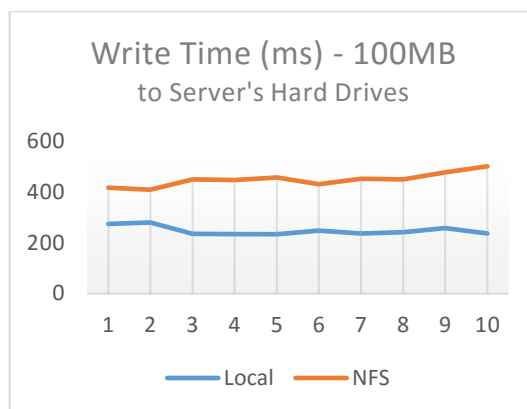


Fig 23. Write Time to Server's Hard Drive – 100MB data size

The Table XV shows the average of 1000 times write testing on 100MB data size to server's hard drive. The using of local drive as virtual machine datastore is 85.50% faster than NFS on writing test with 100MB data size.

TABLE XV
WRITE TIME TO SERVER'S HARD DRIVE – 100MB DATA SIZE

No.	Local Drive	NFS
1	274.3238606	417.0823187
2	279.8248629	409.0026006
3	234.8885168	449.6597092
4	234.4858092	447.0694994
5	233.5891321	456.4545487
6	248.3101872	430.8372765
7	236.5735981	452.3911742
8	241.9112925	449.6875946
9	258.0254186	477.5778424
10	236.4137314	501.0387831
Avg	242.0793223	449.0801347

Figure 24 shows the values of read time (ms) from client to server's hard drives on 100MB data size.

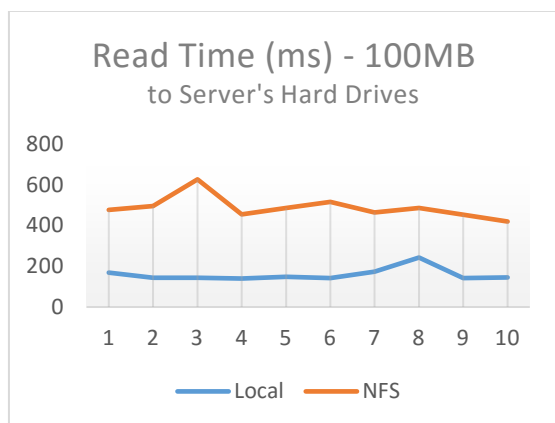


Fig 24. Read Time to Server's Hard Drive – 100MB data size

The Table XVI shows the average of 1000 times read testing on 100MB data size to server's hard drive. The using of local drive as virtual machine datastore is 241.75% faster than NFS on reading test with 100MB data size.

TABLE XVI
READ TIME TO SERVER'S HARD DRIVE – 100MB DATA SIZE

No.	Local Drive	NFS
1	169.3189876	477.4720540
2	143.0246905	495.5157169
3	143.8852889	626.7372177
4	139.7862835	454.7121857
5	148.8176113	485.5323793
6	141.6053269	515.4512313
7	172.7155342	465.0021034
8	242.9524225	486.3185271
9	142.4901530	452.8320866
10	144.6938519	419.7914595
Avg	142.7744716	487.9364962

iii. Round-trip Time from Client to Server

The third testing scenario is measuring RTT value from client to server. The testing was conducted by sending the various size of data packet to server and measure the delay from source to destination and vice versa. When we need to calculate a data and send to be processed on server, we need to know the transmit time needed between client-server, regardless of the time required in the calculations on the server. Figure 25 shows the RTT values between client-server in 4000 times write/read test. The test was conducted when we did testing on schenario 1 (write/read to server's cache memory).

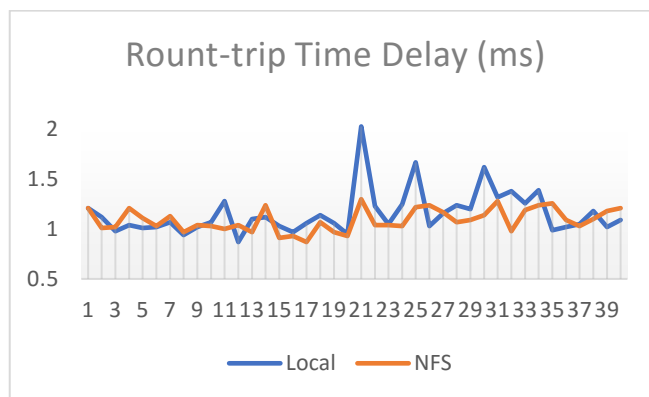


Fig 25. Round-trip Time

The Figure 25 shows the graph from 40 times testing with 100 data packet in a test. The average RTT in 4000 times testing by using local drive is 1.15ms and 1.09ms for using NFS as virtual machine datastore. Therefore, the using of NFS is 6.15% faster than local drive as virtual machine datastore.

VI. CONCLUSION AND FUTURE WORK

This research proposed the using of NFS protocol on NAS storage as vSphere ESXi datastore. The using of NFS can improve manageability on upgrading server's storage space. The test results on write process from client to server's cache memory shows the using of NFS are more stable and faster than local drive as virtual machine datastore on 100KB and 1MB data size, until 10MB and 100MB data size testing that shows the using of local drive become more faster than NFS. The write process to server's cache memory shows contradiction between NFS and local drive as virtual machine datastore on the used of data size. The test on smallest data size (100KB) shows NFS are faster than local drive. However, after the data size is enlarged, the results shows increased speed on the using of local drive and decrease the speed on the use of NFS as virtual machine datastore.

Same as the results of write test, the read test to server's cache memory showed the same pattern. On the testing using a small data shows the use of NFS are faster than local drive. After the data size is enlarged, the results shows the using of local drive becomes better than NFS. Therefore, the use of NFS as virtual machine datastore is suitable for data communication on the cache memory that uses a small data size. The local drive is suitable for data communication on the cache memory that uses a bigger data size.

The test results on write/read speed from client to server's hard drive shows the using of local drive at host hardware as virtual machine datastore are faster than using NFS. The write speed by using local drive as datastore approximately 70% faster than using NFS. The read speed by using local drive approximately 250% faster than using NFS.

The RTT test shows the using of local drive and NFS approx showed almost the same performance. The using of

NFS as virtual machine datastore is just 6.15% faster than local drive after 4000 times testing with various data size.

There are still needed to measure the using of NFS as virtual machine datastore in more diverse data size to get more accurate results related to performance compared to using local drive at host hardware. Also considering the security aspects by using NFS as virtual machine datastore, for the ease of data tapping over the network when the write/read is in process. The security enhancement in the use of NFS as datastore can be done by configuring security protocol for authentication and encrypted data communication. However, there are still need further testing to measure the impact of security methods which will be used related to the data communication performance.

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